

## Lecture 15 - Nov. 3

### Syntactic Analysis

***Identifying Derivations: TDP vs. BUP***  
***Top-Down Parsing: Algorithm***  
***Left-Recursive CFG***

## Announcements

- **Assignment 2** released
  
- **Project Milestone 1** next week
  - + Source project due at 11:59 PM on Tuesday
  - + A simple readme.txt file explains how to run your tool: e.g.,  
    java -jar compiler.jar prog.txt test.txt  
    (and where to find the output HTML file)
  - + Example files you supplied are supposed to work automatically
  - + Jackie will share his screen to build, run, and explore your code.
  
- **Visitor Pattern** source code: Type Casting

# Project: Milestone 1

Milestone 1: Show 3 Example Runs

[ 1% ]

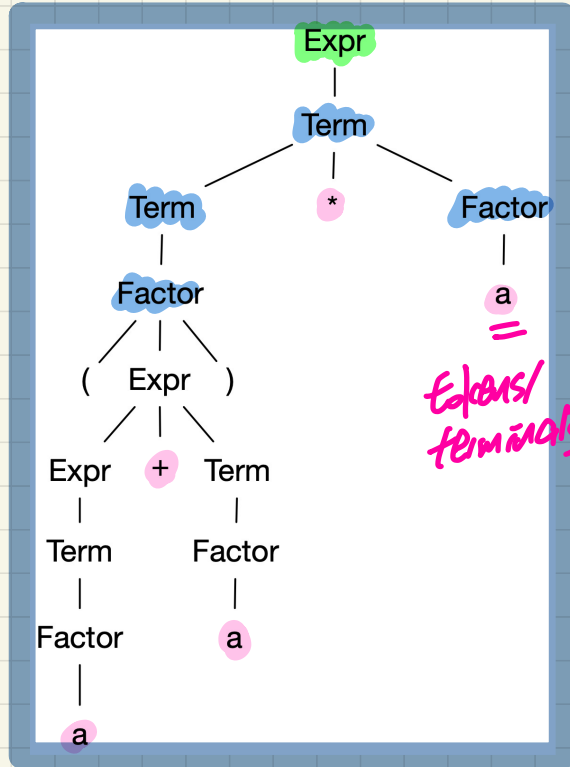
- On the **week of November 7** (about 3 weeks after the project is released), your team is required to meet with Jackie and demonstrate:
  - 3 example runs of your compiler. Each example run consists of the input files and the automatically generated output files.
  - Your example input files should cover (some of the) basic programming features (written in syntax of your own design):
    - ◇ class/module declarations
    - ◇ variable declarations
    - ◇ variable assignments
    - ◇ variable references (i.e., referring to declared variables in expressions)
    - ◇ arithmetic, relational, and logical expressions
    - ◇ conditionals
  - The corresponding produced outputs should cover **at least one control-flow** coverage criterion and **at least one data-flow** coverage criterion.
- **In this meeting, Jackie may suggest specific tasks that your team should complete and will be included in the evaluation of Milestone 2.**

# Discovering Derivations

## Input Grammar G

$Expr$	$\rightarrow$	$Expr + Term$
		$Term$
$Term$	$\rightarrow$	$Term * Factor$
		$Factor$
$Factor$	$\rightarrow$	$(Expr)$
		$a$

AST:  $(a + a) * a$

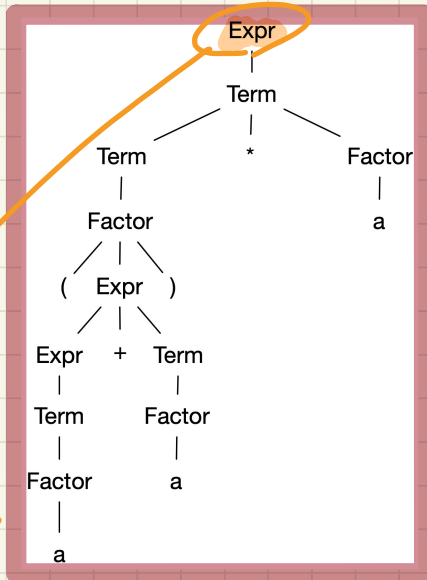


# Discovering Derivations: Top-Down vs. Bottom-Up

## Input Grammar G

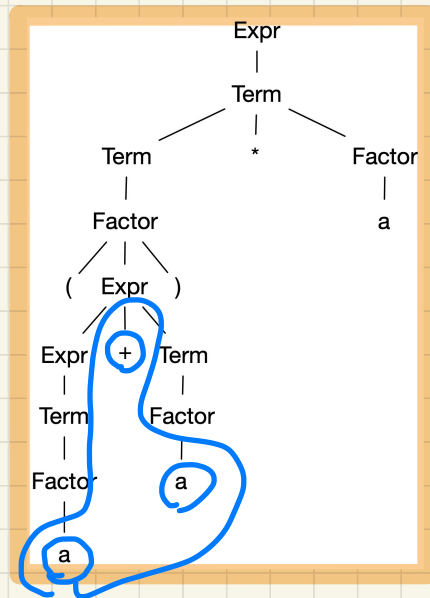
<i>Expr</i>	→	<i>Expr</i> + <i>Term</i>
		<i>Term</i>
<i>Term</i>	→	<i>Term</i> * <i>Factor</i>
		<i>Factor</i>
<i>Factor</i>	→	( <i>Expr</i> )
		a

TDP: (a + a) \* a



possibility of backtrack.

BUP: (a + a) \* a



# Top-Down Parsing: Algorithm

**ALGORITHM:** *TDParse*

**INPUT:** CFG  $G = (V, \Sigma, R, S)$  *string.*

**OUTPUT:** Root of a Parse Tree or **Syntax Error**

**PROCEDURE:**

```

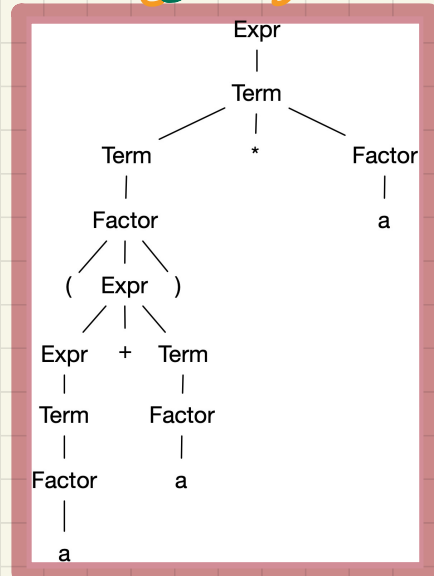
root := a new node for the start symbol S
focus := root
initialize an empty stack trace
trace.push(null)
word := NextWord()
while (true):
  if focus ∈ V then
    if ∃ unvisited rule focus → β1β2...βn ∈ R then
      create β1, β2...βn as children of focus
      trace.push(βnβn-1...β2)
      focus := β1
    else
      if focus = S then report syntax error
      else backtrack
  ✓ elseif word matches focus then
    word := NextWord()
  ✓ focus := trace.pop()
  elseif word = EOF ∧ focus = null then return root
  else backtrack
    
```

*a linear input token seq. into a non-linear AST.*

## Input Grammar G

<b>Expr</b>	→	Expr + Term
		Term
<b>Term</b>	→	Term * Factor
		Factor
<b>Factor</b>	→	(Expr)
		a

**TDP:** (a + a) \* a



**backtrack**  $\triangleq$  pop *focus.siblings*; *focus* := *focus.parent*; *focus.resetChildren*

*SUCCESS*

Expr  $\rightarrow$  Expr + Term  
 $\beta_1 \beta_2 \beta_3$

# Top-Down Parsing: Discovering Leftmost Derivations (1)

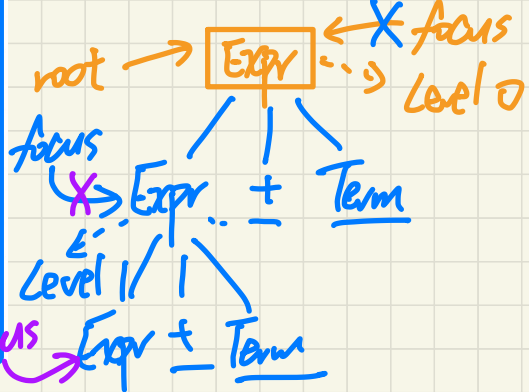
```

ALGORITHM: TDParse
INPUT: CFG  $G = (V, \Sigma, R, S)$ 
OUTPUT: Root of a Parse Tree or Syntax Error
PROCEDURE:
  root := a new node for the start symbol S
  focus := root
  initialize an empty stack trace
  trace.push(null)
  word := NextWord()
  while (true):
    if focus  $\in V$  then
      if  $\exists$  unvisited rule focus  $\rightarrow \beta_1 \beta_2 \dots \beta_n \in R$  then
        create  $\beta_1, \beta_2 \dots \beta_n$  as children of focus
        trace.push( $\beta_n \beta_{n-1} \dots \beta_2$ )
        focus :=  $\beta_1$ 
      else
        if focus = S then report syntax error
        else backtrack
    elseif word matches focus then
      word := NextWord()
      focus := trace.pop()
    elseif word = EOF  $\wedge$  focus = null then return root
    else backtrack
  
```

Parse: **a + a \* a**

Expr	$\rightarrow$	Expr + Term
		Term
Term	$\rightarrow$	Term * Factor
		Factor
Factor	$\rightarrow$	(Expr)
		a

left-recursive



attempts to find LMD

not-terminating

**backtrack**  $\triangleq$  pop focus.siblings; focus := focus.parent; focus.resetChildren

word: "a"

# Left-Recursions (LRs): Direct vs. Indirect

## Direct Left-Recursions:

<i>Expr</i>	→	<i>Expr</i> + <i>Term</i>
		<i>Term</i>
<i>Term</i>	→	<i>Term</i> * <i>Factor</i>
		<i>Factor</i>
<i>Factor</i>	→	( <i>Expr</i> )
		a

<i>Expr</i>	→	<i>Expr</i> + <i>Term</i>
		<i>Expr</i> - <i>Term</i>
		<i>Term</i>
<i>Term</i>	→	<i>Term</i> * <i>Factor</i>
		<i>Term</i> / <i>Factor</i>
		<i>Factor</i>

## Indirect Left-Recursions:

<i>A</i>	→	<i>Br</i>
<i>B</i>	→	<i>Cd</i>
<i>C</i>	→	<i>At</i>

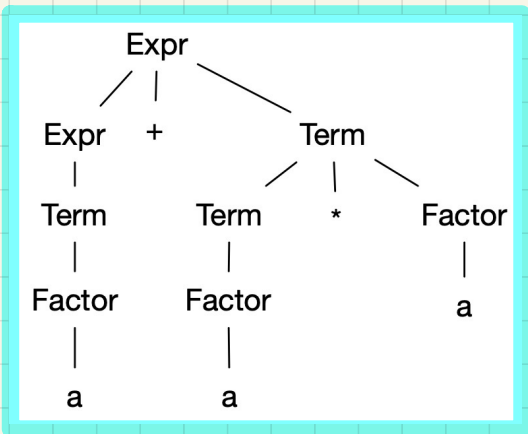
<i>A</i>	→	<i>Ba</i>		<i>b</i>
<i>B</i>	→	<i>Cd</i>		<i>e</i>
<i>C</i>	→	<i>Df</i>		<i>g</i>
<i>D</i>	→	<i>f</i>		<i>Aa</i>   <i>Cg</i>



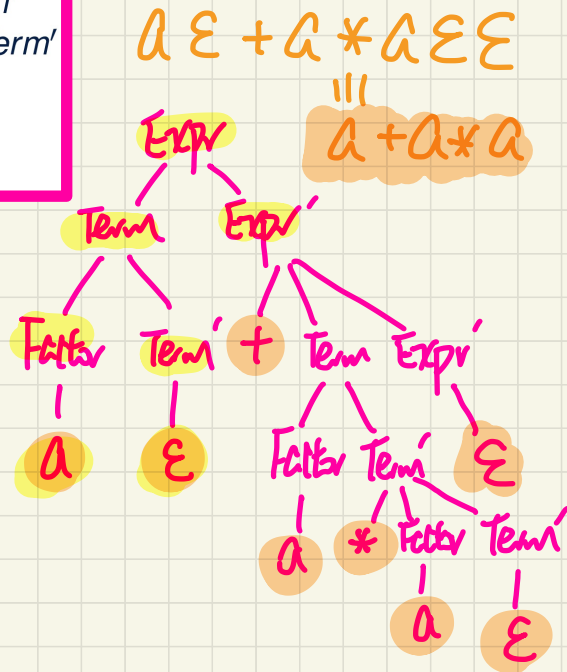
# CFGs: Left-Recursive vs. Right-Recursive

**Example:**  $a + a * a$

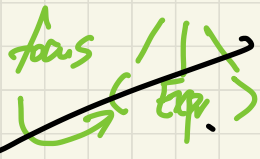
## CFG with Left Recursions

$$\begin{aligned} \text{Expr} &\rightarrow \text{Expr} + \text{Term} \\ &\quad | \quad \text{Term} \\ \text{Term} &\rightarrow \text{Term} * \text{Factor} \\ &\quad | \quad \text{Factor} \\ \text{Factor} &\rightarrow (\text{Expr}) \\ &\quad | \quad a \end{aligned}$$


## CFG with Right Recursions

$$\begin{aligned} \text{Expr} &\rightarrow \text{Term} \text{Expr}' \\ \text{Expr}' &\rightarrow + \text{Term} \text{Expr}' \\ &\quad | \quad \epsilon \\ \text{Term} &\rightarrow \text{Factor} \text{Term}' \\ \text{Term}' &\rightarrow * \text{Factor} \text{Term}' \\ &\quad | \quad \epsilon \\ \text{Factor} &\rightarrow (\text{Expr}) \\ &\quad | \quad a \end{aligned}$$


# Top-Down Parsing: Discovering Leftmost Derivations (2)

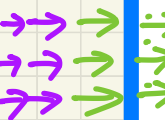
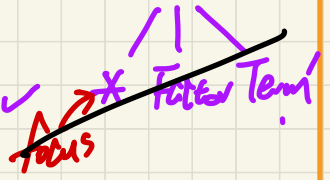


```

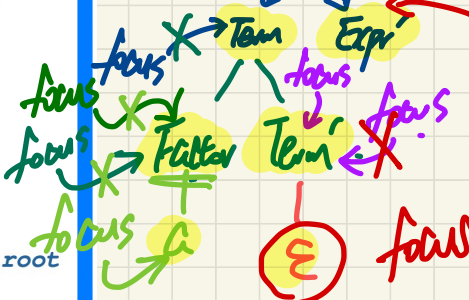
ALGORITHM: TDParse
INPUT: CFG G = (V, Σ, R, S)
OUTPUT: Root of a Parse Tree or Syntax Error
PROCEDURE:
  root := a new node for the start symbol S
  focus := root
  initialize an empty stack trace
  trace.push(null)
  word := NextWord()
  while (true):
    if focus ∈ V then
      if ∃ unvisited rule focus → β1β2...βn ∈ R then
        create β1, β2...βn as children of focus
        trace.push(βnβn-1...β2)
        focus := β1
      else
        if focus = S then report syntax error
        else backtrack
    elseif word matches focus then
      word := NextWord()
      focus := trace.pop()
    elseif word = EOF ∧ focus = null then return root
    else backtrack
  
```

Parse: a + a \* a

Expr	→	Term Expr'	✓
Expr'	→	+ Term Expr'	
			ε
Term	→	Factor Term'	✓
Term'	→	* Factor Term'	
			ε
Factor	→	(Expr)	✓
			a



Next: EVERCRP



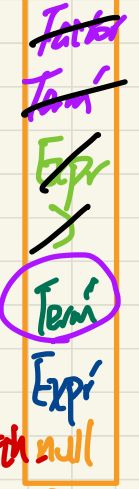
to clarify on Tuesday!



doesn't match "+" but ε has no effect on concat next symbol with null may trace

backtrack ≜ pop focus.siblings; focus := focus.parent; focus.resetChildren

word: "a" "a" "+"



# Top-Down Parsing: Discovering **Leftmost** Derivations (3)

```
ALGORITHM: TDParse
INPUT: CFG  $G = (V, \Sigma, R, S)$ 
OUTPUT: Root of a Parse Tree or Syntax Error
PROCEDURE:
  root := a new node for the start symbol  $S$ 
  focus := root
  initialize an empty stack trace
  trace.push(null)
  word := NextWord()
  while (true):
    if focus  $\in V$  then
      if  $\exists$  unvisited rule  $focus \rightarrow \beta_1\beta_2\dots\beta_n \in R$  then
        create  $\beta_1, \beta_2, \dots, \beta_n$  as children of focus
        trace.push( $\beta_n\beta_{n-1}\dots\beta_2$ )
        focus :=  $\beta_1$ 
      else
        if focus =  $S$  then report syntax error
        else backtrack
    elseif word matches focus then
      word := NextWord()
      focus := trace.pop()
    elseif word = EOF  $\wedge$  focus = null then return root
    else backtrack
```

**Parse:**  $(a + a) * a$

<i>Expr</i>	$\rightarrow$	<i>Term</i>	<i>Expr'</i>
<i>Expr'</i>	$\rightarrow$	$+$	<i>Term</i> <i>Expr'</i>
			$\epsilon$
<i>Term</i>	$\rightarrow$	<i>Factor</i>	<i>Term'</i>
<i>Term'</i>	$\rightarrow$	$*$	<i>Factor</i> <i>Term'</i>
			$\epsilon$
<i>Factor</i>	$\rightarrow$	$($	<i>Expr</i> )
			$a$

**EXERCISE**

**backtrack**  $\triangleq$  pop *focus.siblings*; *focus* := *focus.parent*; *focus.resetChildren*